UNSYMMETRIC ORDERING USING A CONSTRAINED MARKOWITZ SCHEME

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Abstract. We consider the $LU$ factorization of unsymmetric sparse matrices using a three-phase approach (analysis, factorization and triangular solution). Usually the analysis phase first determines a set of potentially good pivot and then orders this set of pivots to decrease the fill-in in the factors. In this paper, we present a preprocessing algorithm that simultaneously achieves the objectives of selecting numerically good pivots and preserving the sparsity. We describe the algorithmic properties and difficulties in implementation. By mixing the two objectives we show that we can reduce the amount of fill in the factors and reduce the number of numerical problems during factorization. On a set of large unsymmetric real problems, we obtain the average gains of 14% in the factorization time, of 12% in the size of the $LU$ factors, and of 21% in the number of operations performed in the factorization phase.

1. Introduction. We consider the $LU$ factorization of a unsymmetric sparse matrix $A$ based on a three-phase approach. It includes an analysis phase, a numerical factorization phase, and a triangular solution phase. In this class of methods, without loss of generality, we will consider the multifrontal and the supernodal algorithms (see for example [2, 8, 12, 17, 19]). The analysis phase transforms $A$ into $\tilde{A}$ with better properties for sparse factorization. It exploits the structural information to reduce the number of fill-ins in the $LU$ factors and exploits the numerical information to reduce the number of numerical pivotings needed during factorization. Two separate treatments are commonly used in sequence for these two objectives. Firstly, scaling and maximum transversal algorithms are used to transform $A$ into $A_1$ with large entries in magnitude on the diagonal. Secondly, a symmetric fill-reducing ordering, which preserves the large diagonal, is used to permute $A_1$ into $A_2$ so that the factors of $A_2$ are sparser than those of $A_1$. Note that during factorization, numerical instabilities can still occur and will be handled either by partial pivoting resulting in extra fill-ins in the factored matrices or by static pivoting resulting in a potentially less accurate factorization. This approach has two drawbacks:

- The numerical treatment forces the fill-reducing ordering to restrict pivot selection to the diagonal of $A_1$, and so to compute a symmetric permutation,
- The ordering phase does not have numerical information to select pivots.

To reduce the fill-in in the factors and improve the numerical quality of the preselected pivots, we describe in this paper, a family of orderings that can select off-diagonal pivots based on a combination of structural and numerical criteria. Based on a numerical preprocessing of the matrix we build a set of numerically acceptable pivots, referred to as matrix $C$, that may contain off-diagonal entries. We then compute an unsymmetric ordering taking into account both the structure of $A$ and the numerical

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information in C. The C matrix serves as a constraint matrix for the pivot selection. The new algorithm is referred to as constrained Markovs with local symmetrization (or OMLS). In this context, the use of local symmetrization introduced in [4] will be explained and justified in Section 3.2.

We will focus on two state-of-the-art direct solvers representative of this class: the multifrontal code MA41JNS [2, 5] and the supernodal code SuperLU_DIST [27, 28]. For the MA41JNS multifrontal code, standard partial pivoting with a threshold value is applied to locally select numerically stable pivots within a so-called frontal matrix. It is possible that some variables cannot be locally eliminated. Postponing the elimination of a pivot will then result in an increase in the size of the LU factors estimated during analysis and in an increase in the number of operations. In the case of the supernodal code, SuperLU_DIST, the distributed memory version uses a right-looking formulation which, having computed the factorization of a block of columns, then immediately sends the data to update the block columns in the trailing submatrix. A static pivoting strategy is used and we keep the pivotal sequence chosen in the analysis. The result is that iterative refinement may be needed to improve the solution.

In practice, it has been observed in [3] that setting large entries on the diagonal, based on [14], one can significantly reduce the number of numerical problems for both solvers. Nevertheless, the static approach can still fail on some problems and relies more on iterative refinement in the solution phase. In this context, the problem of finding a good ordering for preserving the diagonal entries (symmetric ordering) and taking into account the asymmetry of the matrix was studied in [4]. In this paper the work described in [4] (referred to as DMLS) has been extended and generalized as follows. Firstly, we do not limit our choice of pivots to a transversal of the original matrix. Secondly, at each step of elimination, the pivot is chosen within a constraint matrix C, which is updated at each step. Thirdly, we may also consider numerical values and numerical errors during the pivot selection. To do so, we compute an incomplete factorization of the constraint matrix. Finally, because of all the new concepts introduced in this approach, we had to introduce new algorithms, revisit the existing ones and introduce new data structures to implement our algorithms.

The rest of the paper is organized as follows. Section 2 introduces the main components of our algorithm. Section 3 defines the graph-theoretic notations and describe the use of local symmetrization in our context. Section 4 describes the algorithmic contributions of the proposed OMLS method. A full detailed presentation of our implementation is given in [32]. Section 5 analyses the results of the newly implemented OMLS algorithm when applied to real-life unsymmetric test cases.

2. Components of our unsymmetric ordering. Given a matrix A, our unsymmetric ordering consists of two main steps:

- **Step 1**: Based on a numerical pre-treatment of the matrix A, we extract a set of numerically acceptable pivots, referred to as the constraint matrix C. We have Pattern(C) ⊂ Pattern(A), and if c_{ij} ≠ 0 then c_{ij} = a_{ij}.
- **Step 2**: Constrained unsymmetric ordering; the constraint matrix is used at each step of the symbolic Gaussian elimination to control the set of eligible pivots (possibly with respect to both numerical and structural criteria).

Before describing more precisely these 2 steps, we introduce definitions and notations that will be used to describe our algorithms.

Let M = (m_{ij}) be a matrix of order n. If M can be permuted to have nonzeros on the diagonal then M is structurally nonsingular. Let \( G_M = (V_r, V_c, E) \) be the
bipartite graph associated with the matrix $M$. $V_r$ is the set of row vertices and $V_c$ is the set of column vertices. Let $(i, j) \in V_r \times V_c$ then $(i, j) \in E$ if and only if $m_{ij} \neq 0$. A matching is a subset of edges $M \subset E$ such that for all vertices $v \in V$, at most one edge of $M$ is incident on $v$. If $M$ is structurally nonsingular, then there exists a matching $M$ with $n$ edges and $M$ is said to be a perfect matching. We will say also that $M$ is a perfect transversal.

For the sake of clarity, in the remainder of this paper we assume that $A$ is structurally nonsingular. The adaptation of our algorithms to structurally singular matrices is straightforward but would have severely over-complicate our notations and our comments.

2.1. Step 1: Numerical preprocessing. The objective of this preprocessing step is to extract the most significant (structurally and numerically) entries of matrix $A$ and to use them to build the constraint matrix $C$.

Firstly, we scale the matrix $A$ with the diagonal matrices $D_r$ and $D_c$, resulting in $A \leftarrow D_r AD_c$. Secondly, a constraint matrix $C$ can be constructed from $A$ such that $\text{Pattern}(C) \subset \text{Pattern}(A)$ and $C$ satisfies certain numerical and/or structural properties. Since the entries in $C$ correspond to the potential pivots for the subsequent step, we only keep a subset of bounded size (typically less than $3n$) of the largest entries in the scaled matrix. Furthermore, we want $C$ to be structurally nonsingular and thus we add entries from $A$ to guarantee that $C$ includes a perfect transversal $M$. Note that $M$ is also a perfect transversal of $A$.

2.2. Step 2: Constrained unsymmetric ordering. Let $A^1 = A$ be the original matrix of order $n$ and $A^k$ be the reduced matrix after eliminating the first $k - 1$ pivots (not necessarily on the diagonal). Let $C^1 = C$ be such that $\text{Pattern}(C^1) \subset \text{Pattern}(A^1)$. At each step $k$, a pivot $p^k$ such that $p^k \in \text{Pattern}(C^k)$ is selected. This selection may combine structural heuristics based on the structure of $A^k$ (e.g., approximate Markowitz count, approximate minimum fill, etc.) and numerical heuristics carried by the $C^k$ matrix. Matrix $A^k$ is updated (remove row and column of the pivot and add fill-ins in the Schur complement). Matrix $C^k$ is updated such that $C^{k+1}$ remains structurally nonsingular and $\text{Pattern}(C^{k+1})$ is included in the pattern of $C^k$, where $C^k$ is defined as the reduced matrix after the elimination of pivot $p^k$ in $C^k$. Note that $C^k$ includes the fill-in resulting from the elimination of $p^k$ in $A^k$. This implies that $\text{Pattern}(C^{k+1}) \subset \text{Pattern}(A^{k+1})$. To keep $C^k$ structurally nonsingular, a perfect matching in $C^k$ is maintained at each step. When there is no ambiguity, we will omit the superscript $k$ from the matrix notations.

The following two considerations influence the update that will be performed on $C$:

- Which metric do we use to select a pivot?
- Which entries and/or values are added/updated in $C$ at each step of the elimination?

Note that if we consider the magnitude of the $C$’s entries to select a pivot, both the pattern of $C$ and the numerical values need be stored and updated. Furthermore, the structural metric of each entry $(i, j)$ in $C$ should carry information on the reduced matrix associated with the complete matrix $A$.

The ordering algorithm also depends on how $C$ is updated at each step. As mentioned before in the description of Step 2, we want at each step to guarantee that:

\[(2.1) \quad C \text{ must remain structurally nonsingular,}\]
3. Notations and definitions. Before giving the algorithmic details of the proposed OMLS method, we introduce the graph structures and notations that will be used in this paper. The structure of an unsymmetric matrix can be represented as a bipartite graph, and Gaussian elimination can be efficiently modeled by a bipartite quotient graph [31]. We first describe the main properties of bipartite graphs and bipartite quotient graphs and their relationship with Gaussian elimination. We then introduce the notations that will be used to describe our algorithms and define local symmetrization [4], a technique that simplifies the bipartite quotient graph implementation. Note that we use calligraphic letter for notations related to quotient graphs.

3.1. Bipartite graph. Let \( C = (c_{ij}) \) be a matrix and \( G_C = (V_r, V_c, E) \) be its associated bipartite graph. Let \( R_i \) denotes the structure of row \( i \), i.e., \( R_i = \{ j \in V_c \; s.t. \; (i,j) \in E \} \). Let \( C_j \) denote the structure of column \( j \), i.e., \( C_j = \{ i \in V_r \; s.t. \; (i,j) \in E \} \).

In Gaussian elimination, when a pivot \((i,j)\) is eliminated, a new matrix, referred to as the reduced matrix \( \tilde{C} \) is computed. \( \tilde{C} \) is obtained from \( C \) by removing row \( i \) and column \( j \) and by adding the Schur complement entries. In terms of graph manipulations, this elimination adds edges in the bipartite graph of \( C \) to connect all the rows adjacent to \( j \) to all the columns adjacent to \( i \). This set of connected rows and columns is referred to as a bi-clique.

The symbolic factorization of \( C \) is done by building \( C^k \) for \( k = 1 \) to \( n \), with \( C^1 = C \). After eliminating the \( k^{th} \) pivot \((row_p, col_p)\) in \( C^k \), we compute \( C^{k+1} = \tilde{C}^k \).

More formally, the changes of the bipartite graph from \( G_{C^k} \) to \( G_{C^{k+1}} \) are

\[
R_i^{k+1} = (R_i^k \cup R_{row_p}^k) \setminus \{col_p\} \quad \text{for} \quad i \in C_{col_p}^k,
\]

\[
C_j^{k+1} = (C_j^k \cup C_{col_p}^k) \setminus \{row_p\} \quad \text{for} \quad j \in R_{row_p}^k,
\]

and remove \( row_p \) and \( col_p \) from the vertex sets of \( G_{C^{k+1}} \). When it is clear from the context, we will use the notation \( L_p \) in place of \( R_{row_p}^p \) and \( U_p \) in place of \( C_{col_p}^p \).

3.2. Bipartite quotient graph. In the previous section we have shown that, to update the bipartite graph we must add, at each elimination step, the entries to the Schur complement matrix which may be costly to update and to store. It has been shown that the quotient graphs can be used to efficiently model the factorization of symmetric matrices [18, 22]. The main idea is to use a compact representation of the cliques associated with the eliminated vertices. This concept can be extended (see [31]) to model the LU factorization. In this case, a bipartite quotient graph can be used to represent the edges in a bi-clique. It has then been shown in [31] that doing so the elimination can be modeled in space bounded by the size of the original matrix \( A \).

In this section, we first explain why the quotient graph model leads to more complex algorithms on unsymmetric matrices than on symmetric matrices. We then briefly define element absorption and explain the use of local symmetrization to reduce the quotient graph complexity. Finally we introduce the notations that will be used to describe our algorithms.

Let \( G^k \) be the bipartite quotient graph used to represent the structure of the reduced submatrix \( A^k \) after \( k \) steps of elimination. Initially the bipartite quotient graph \( G^1 \) is identical to the bipartite graph \( G^1 \). At step \( k \) of Gaussian elimination,
any eliminated pivot \( e = (r_e, c_e) \) will be referred to as a coupled row and column element. All the row and column vertices that are not coupled elements are referred to as the row and column variables of \( G^k \). Both row and column vertices of the graph are thus partitioned into two sets composed of variables (uneliminated vertices) and elements (eliminated vertices). We then define \( G^k = (V_r \cup V_c, V_c \cup \overline{V}_c, \mathcal{E} \cup \overline{\mathcal{E}}) \). The vertices in \( V_r \) (respectively \( V_c \)) correspond to the row (respectively column) variables. The vertices in \( \overline{V}_r \) (respectively \( \overline{V}_c \)) correspond to the row (respectively column) elements. The edge set \( \mathcal{E} \) is such that \( \mathcal{E} \subseteq (V_r \times V_c) \) whereas \( \overline{\mathcal{E}} \) is such that \( \overline{\mathcal{E}} \subseteq (V_r \times \overline{V}_c) \cup (\overline{V}_r \times V_c) \cup (\overline{V}_r \times \overline{V}_c) \). With our definitions \((r, c)\) is a nonzero entry in the reduced matrix at step \( k \) if and only if there exists a path joining \( r \) and \( c \) which only visits the elements and for which all the edges in the even positions correspond to already eliminated pivots. In other words, the structure of a row \( i \) at step \( k \) is the set of reachable columns \( j \) through all the paths of the form \( i \to c_{e_1} \to r_{e_1} \cdots \to c_{e_l} \to r_{e_l} \to \cdots \to c_{e_l} \to \cdots \to c_{e_1} \to i \) where \( e_i = (r_{e_i}, c_{e_i}) \), \( 1 \leq i \leq l \) are coupled elements. Similarly, the structure of a column \( j \) at step \( k \) is the set of reachable rows \( i \) through all the paths of the form \( j \to r_{e_1} \to c_{e_1} \cdots \to r_{e_l} \to c_{e_l} \to \cdots \to c_{e_1} \to i \). This process may involve paths of arbitrary length in \( G^k \) and in particular through more than one coupled element.

For example in Figure 3.1, we assume that the entry \((\text{row}_p, c_{e_2})\) corresponds to the fill-in due to the elimination of element \( e_1 \) and is implicitly represented by an edge between variable \( \text{row}_p \) and element \( c_{e_1} \) in the quotient graph. In this case we know that the row structure of \( \text{row}_p \) contains the row structure of \( e_2 \) because of the path \( \text{row}_p \to c_{e_1} \to r_{e_1} \to c_{e_2} \to r_{e_2} \).

In the context of sparse Cholesky factorization, an undirected quotient graph (the row and column vertices are merged) is preferred and commonly used to compute an ordering for symmetric matrices (e.g., Multiple Minimum Degree [29] and Approximate Minimum Degree [1]). The structure of the factors can be computed following the paths of length at most two in this quotient graph. There are no edges between the elements.

In the unsymmetric case, when a pivot \( p = (\text{row}_p, \text{col}_p) \) is selected, if there exists a cycle of the form \( \text{row}_p \to c_{e_1} \to r_{e_1} \cdots \to c_{e_l} \to r_{e_l} \to \text{col}_p \to \text{row}_p \) (referred to as a strongly connected component in [31]) then, except for \( \text{row}_p \) and \( \text{col}_p \), the row and column elements in the cycle are not needed any more to retrieve the structure of the remaining variables. Firstly, all the row (resp. column) variables adjacent to these column (resp. row) elements will be included in the adjacency of \( \text{col}_p \) (resp. \( \text{row}_p \)), and secondly, the row (resp. column) variables which were adjacent to one of the removed elements will be adjacent to \( \text{col}_p \) (resp. \( \text{row}_p \)). This process will be called element absorption and is illustrated in Figure 3.1.

To avoid long search paths, we decided to relax the element absorption rule. A row (resp. column) element is absorbed by the current row (resp. column) pivot if either it is adjacent to the column (resp. row) pivot or its associated column (resp. row) element is adjacent to the row (resp. column) pivot. This is referred to as local symmetrization in [4]. It implies that the resulting quotient graph \( G^k \) at step \( k \) models only an approximation of the structure of the reduced submatrix. It has been shown in [4] that the exploitation of element absorption combined with local symmetrization results in an in-place algorithm: at each step of the Gaussian elimination, the size of the quotient graph is bounded by the size of \( G^1 \). Note that because of local symmetrization, an approximation of the symbolic factors can be computed following the paths of length at most three of the form \( i \to c_e \to r_e \to j \) where \((r_e, c_e)\) denotes a coupled row and column element.
Fig. 3.1. Illustration of a cycle \((\text{row}_p \rightarrow c_{e_1} \rightarrow r_{e_1} \rightarrow c_{e_2} \rightarrow r_{e_2} \rightarrow \text{col}_p \rightarrow \text{row}_p)\).

To simplify the description of how the bipartite quotient graph is modified at each elimination step, we define \(\overline{V} \subset (\overline{V}_r \times \overline{V}_c)\) to be the set of coupled row and column elements corresponding to already eliminated pivots. Entries of the set \(\overline{V}\) will also be referred to as coupled elements or elements when it is clear from the context. Let \(U_p\) (resp. \(L_p\)) be the row (resp. column) variables adjacent to the row (resp. column) element of a pivot \(p = (\text{row}_p, \text{col}_p)\). Thanks to local symmetrization, the concept of absorption can be extended to coupled elements: an element \(e = (r_e, c_e)\) such that \((\text{row}_p, c_{e}) \in \overline{E}\) or \((r_e, \text{col}_p) \in \overline{E}\) can be absorbed by \(p\) when \(p\) is selected as a pivot. A consequence of this absorption is that our ordering also generates a dependency graph between elements that is in fact a forest. This forest will be fully exploited by the unsymmetrized multifrontal approach [5].

For each row variable \(i \in \mathcal{V}_r\) and column variable \(j \in \mathcal{V}_c\), we define the element lists \(\mathcal{R}_i\) and \(\mathcal{C}_j\) as follows:

\[
\mathcal{R}_i = \{ e = (r_e, c_e) \in \overline{V} \text{ s.t. } (i, c_e) \in \overline{E} \}
\]

and

\[
\mathcal{C}_j = \{ e = (r_e, c_e) \in \overline{V} \text{ s.t. } (r_e, j) \in \overline{E} \}.
\]

Let \(e = (r_e, c_e)\) be an element, if \(e \in \mathcal{R}_i\) then we will say that element \(e\) is adjacent to row variable \(i\). Similarly, if \(e \in \mathcal{C}_j\) we will say that element \(e\) is adjacent to column \(j\).

Using this notation, the adjacency of a row variable \(i\) (resp. column \(j\)) in \(G\) consists of a list of column variables denoted as \(\mathcal{A}_{i*}\) (resp. a list of row variables \(\mathcal{A}_{*j}\)) and a list of elements \(\mathcal{R}_i\) (resp. \(\mathcal{C}_j\)). At the beginning \(\mathcal{R}_i = \mathcal{C}_j = \emptyset\) and \(\mathcal{A}_{i*} = \mathcal{A}_{i*}^{(0)}\) and \(\mathcal{A}_{*j} = \mathcal{A}_{*j}^{(0)}\) corresponds to the original entries of \(A\). Each step of Gaussian elimination involves changes in the sets \(\mathcal{R}_i\) and \(\mathcal{C}_j\) as well as the computation of the structure of a current pivot \(p\) (computation of \(L_p\) and \(U_p\)). The variable lists \(\mathcal{A}_{i*}\) and \(\mathcal{A}_{*j}\) can also be pruned. Indeed, the edges in \(G\) between the variables and the elements implicitly represent the bi-clique of the element and can thus be used to remove the redundant entries in \(\mathcal{A}_{i*}\) and \(\mathcal{A}_{*j}\). This important point will be further discussed in detail in Section 4.3.

When \((i, j) \in \mathcal{V}_r \times \mathcal{V}_c\) is selected as the next pivot we build the element \(p = \)
\((i, j) \in \mathcal{G}\) such that:

\[
U_p = A_{i*} \cup \bigcup_{e \in R_i} U_e \cup \bigcup_{e \in C_j} U_e
\]

and

\[
L_p = A_{*j} \cup \bigcup_{e \in C_j} L_e \cup \bigcup_{e \in R_i} L_e.
\]

The third term in each equation results from local symmetrization and will enable the current pivot to absorb all the elements which it was adjacent to. For example, let us assume that the entry \(p1\) is selected as pivot in Figure 3.2. Since \(col_p\) is adjacent to \(e1\), local symmetrization adds the virtual \(Sp1\) entry so that the row structure of \(p1\) contains \(U_{e1}\). Let \(F_p = C_j \cup R_i\) be the set of elements adjacent to the current pivot. Note that when it is clear from the context, \(F\) will also refer to the elements adjacent to the current variable. The elements in \(F_p\) are absorbed by \(p\) and can be removed from the quotient graph. For each column variable \(j1\) in \(U_p\) (resp. \(i1\) in \(L_p\)) the element \(p\) is added to its adjacency and the elements in \(F_p\) are removed: \(C_{j1} \leftarrow (C_{j1} \setminus F_p) \cup \{p\}\) (resp. \(R_{i1} \leftarrow (R_{i1} \setminus F_p) \cup \{p\}\)). The structure of column \(j1\) of the factors in the reduced matrix is then given by \(A_{ij1} \cup \bigcup_{e \in C_j} L_e\). The structure of row \(i1\) of the factors is \(A_{i*} \cup \bigcup_{e \in R_i} U_e\).

![Diagram](image)

**Fig. 3.2. Influence of local symmetrization on the pivot structure**

Note that, although the above structural changes of the reduced submatrix are correct, they should not be used to compute the structural metrics. Indeed, if \((i1, j1)\) is selected as the next pivot, then the correctly computed structure of the reduced matrix should include the local symmetrization terms (similar to equations (3.1) and (3.2)). In Figure 3.2, we illustrate the effect of local symmetrization on the structure of the selected pivot. Let us consider two candidate pivots belonging to the same row \(row_p\), \(p1 = (row_p, col_{p1})\) and \(p2 = (row_p, col_{p2})\). We assume that all the elements in \(\mathcal{G}\) adjacent to \(p1\) and \(p2\) are indicated in the figure. The structure of \(row_p\) is then given by \(A_{row_p*} \cup U_{e3}\). This however does not give enough information on the structure of \(row_p\) if either \(p1\) or \(p2\) were selected as the next pivot. If \(p1\) were the next pivot then the structure of \(row_p\) would be given by \(U_{p1} = A_{i*} \cup U_{e3} \cup U_{e1}\) because of the locally symmetrized entry \(Sp1\). If \(p2\) were the next pivot then the structure of \(row_p\)
<table>
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<tr>
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<th>Bipartite quotient graph $G^k$</th>
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<tr>
<td>$G_C = (V_C, V_r, E)$</td>
<td>$G^k = (V_r, V_c, V_C, E')$</td>
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<tr>
<td>$R_i$</td>
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<td>elements adjacent to row $i$</td>
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<td>$C_j$</td>
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<tr>
<td>structure of column $j$</td>
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<td>$U_{p}$</td>
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<td>row structure of the $p^{th}$ pivot</td>
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<td>$L_{p}$</td>
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<tr>
<td>column structure of the $p^{th}$ pivot</td>
<td>column structure of the $p^{th}$ pivot</td>
</tr>
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| $F$ | elements that are adjacent to row $i$ or column $j$ for a pivot $(i, j)$ ($F = R_i \cup C_j$) |

Table 4.1

Notations used for bipartite graph and bipartite quotient graph.

would be given by $U_{p2} = A_{p2} \cup U_{c3} \cup U_{c2}$ because of the locally symmetrized entry $S_{p2}$. This shows that, even if we cannot anticipate the effect of local symmetrization on the quotient graph $G$ before the pivot selection, we should anticipate its effect on the metrics used to select the best pivot between $p_1$ and $p_2$.

4. CMLS algorithm. In this section, we describe the main features and properties of the CMLS algorithm. At each step of the algorithm, we need to know the exact structure and the metric of each nonzero entry in $C$. It is thus natural to use a bipartite graph (with possibly weighted edges) for $C$. Each edge corresponding to a nonzero entry may have one or more weights (for example, a numerical value and a structural metric) that will be used to select a pivot. On the other hand, in order to have a fast computation of the structural metrics based on the pattern of $A$ and to have an in-place algorithm, $A$ is represented by its quotient graph and local symmetrization is employed. The notations used to represent the two graphs at each step of the algorithm are summarized in Table 4.1.

In Section 4.1, we first describe the pivot selection algorithms. Updating the graph $G_C$ and $G$ associated with $C$ and $A$ respectively is discussed in Sections 4.2 and 4.3. In Section 4.4 we describe how to compute, at each step $k$ and for each entry in the constraint matrix $C^k$, the structural metric relative to $G^k$. Section 4.5 finally explains how supernovas are defined and used in our context.

4.1. Pivot selection. At each step, the pivot with minimal metric is selected. The choice of a metric implies the underlying algorithmic strategy. We say that we use structural strategies in our algorithms when the entries are selected with respect to only the structural metrics whereas the hybrid strategies correspond to the combination of structural and numerical metrics.

Using doubly-linked lists and a hash function based on the value of the structural metric provides a direct access to the entry in $C$ that has the minimum structural metric in $G_C$.

Furthermore, as it is often the case in sparse matrix factorization, we may want to preserve the sparsity of the factors while controlling the growth in the factors. Numerical thresholds are introduced to give freedom for the pivot selection to balance numerical precision with sparsity preservation. For each entry $(i, j) \in C^k$ we define its numerical metric: $v(i, j) = |c_{ij}|/||c_j||_{\infty}$. A pivot in position $(i, j)$ is said to be numerically acceptable (or acceptable) according to a threshold $\tau$ if and only if $|c_{ij}| \geq \tau \times ||c_j||_{\infty}$ where $\tau \in [0, 1]$. 
To reduce the complexity of the algorithms, it is also common to limit the pivot search to a set of candidate pivots. For example in [11] the authors proposed to visit the entries of a fixed number of columns using the Zlatev-style search [34]. We use a slightly different algorithm to limit our search. At each step of the ordering, we look for the best entry \( p = (\text{row}_p, \text{col}_p) \) within a subset (say \( S \)) of the entries in the bipartite graph \( G_C \). The subset \( S \) is defined by two threshold parameters \( MS > 0 \) and \( NCOL \geq 0 \) as follows. Firstly, the \( MS \) entries with the smallest structural metric \( m_0 \) are added to \( S \). Secondly, those \( MS \) entries may belong to several different columns. We then add in \( S \) all the other nonzero entries of those columns, but restricted to at most the first \( NCOL \) columns. The set \( S \) is thus composed of a first set of \( MS \) entries, the so-called \( MS \)-set, and a second set, the so-called \( NCOL \)-set = \( S \setminus MS \)-set. This implies that if \( NCOL > 0 \), the \( NCOL \)-set will guarantee that the numerically best pivots in the first \( NCOL \) columns will belong to \( S \). As for the structural strategies, it is straightforward to access entries in the \( MS \)-set and their corresponding columns using the doubly-linked lists and a hash function.

We now explain how we select the entry of minimum structural metric in \( S \) among the numerically acceptable pivots. We first visit the \( MS \)-set sorted in increasing order of the structural metric \( m_0 \). The first numerically acceptable entry found corresponds to the minimum with respect to our hybrid strategy and we stop the search. Otherwise, none of the values in the \( MS \)-set entries is numerically acceptable. However, if \( NCOL > 0 \) then we are sure that at least \( NCOL \) entries will be numerically acceptable since \( \tau \leq 1 \). In this case, however, all the entries of \( S \) need to be considered to obtain the minimum on \( S \) according to our hybrid strategy. Finally if \( NCOL = 0 \) and none of the entries in the \( MS \)-set is numerically acceptable then the first entry of the \( MS \)-set is selected even if it is not numerically acceptable.

4.2. Update of the bipartite graph \( G_C \). A bipartite graph is used to represent \( C \). At each step \( k \), we need to add new entries in \( G_{C^{k+1}} \) corresponding to the fill-ins in \( C^{k+1} \). Since \( G_C \) holds the set of candidate pivots, we need to guarantee that the Properties (2.1) and (2.2) hold.

Let \( \mathcal{M} \) be a matching in \( C^k \). The following two extreme strategies preserves these two properties:

- \textbf{MATCHUPDATE} will refer to the strategy that performs incomplete Gaussian elimination to only preserve the perfect matching property (2.1). Let \( p = (\text{row}_p, \text{col}_p) \) be the current pivot. Let \((\text{row}_p, \text{match} \text{.col})\) and \((\text{match} \text{.row}, \text{col}_p)\) be the matched entries of \( C \) in row \( \text{row}_p \) and column \( \text{col}_p \), respectively. That is, \((\text{row}_p, \text{match} \text{.col}) \in \mathcal{M} \) and \((\text{match} \text{.row}, \text{col}_p) \in \mathcal{M} \). If these entries are the same (i.e., \((\text{row}_p, \text{col}_p) \) is a matched entry), nothing needs to be done to maintain Property (2.1). Otherwise entry \((\text{match} \text{.row}, \text{match} \text{.col})\) is added to maintain Property (2.1). Note that this entry corresponds to an entry in \( \text{Pattern}(C^k) \), so that Property (2.2) remains true.

- \textbf{TOTALUPDATE} will refer to the strategy which performs all the updates in \( C \) (i.e., \( C^{k+1} = C^k \)).

In practice, a mixed strategy, exploiting both MATCHUPDATE and TOTALUPDATE will be used for the experiments. At each step, TOTALUPDATE strategy is used except when we want to limit the memory or the time complexity, then MATCHUPDATE strategy is used. See [32] for further details.
4.3. Update of the bipartite quotient graph $G$. In Algorithm 4.1 we describe how the bipartite quotient graph associated with the reduced matrix is updated.

**Algorithm 4.1** QMLS update of the bipartite quotient graph $G^k$

Let $p = (\text{row}_p, \text{col}_p)$ be the current pivot at step $k$ and $F_p = \mathcal{R}_{\text{row}_p} \cup \mathcal{C}_{\text{col}_p}$.

1. If $U_p \neq \emptyset$ and $L_p \neq \emptyset$, then
   1.1. do
   1.2. for each row $i \in U_p$ do
   1.3. for each column $j \in L_p$ do
   1.4. $A_{ij}$ = $(A_{ij} \setminus U_p) \setminus \{\text{col}_p\}$ /* variable elimination in row direction */
   1.5. end for
   1.6. end for
   1.7. $R_i = (R_i \setminus F_p) \cup \{\text{col}_p\}$
   1.8. $A_i = A_i \setminus \{\text{col}_p\}$
   1.9. end for
   1.10. end for
   1.11. $C_j = (C_j \setminus F_p) \cup \{\text{row}_p\}$
   1.12. $A_{ij} = A_{ij} \setminus \{\text{row}_p\}$
   1.13. end for
   1.14. end for
   1.15. else /* pivot pruning: delete all that is related to $p$, if $U_p = \emptyset$ or $L_p = \emptyset$ */
   2. for each row $i \in L_p$ do
   2.1. $R_i = (R_i \setminus F_p)$
   2.2. $A_i = A_i \setminus \{\text{col}_p\}$
   2.3. end for
   2.4. for each column $j \in U_p$ do
   2.5. $C_j = (C_j \setminus F_p) \cup \{\text{row}_p\}$
   2.6. $A_{ij} = A_{ij} \setminus \{\text{row}_p\}$
   2.7. end for
   2.8. end if

The “if” block of code shows how the elements and variables are pruned. The element pruning performed at lines 2 and 4 includes pruning due to local symmetrization. The variable pruning performed at lines 1 and 3 removes the intersection of the adjacency structures. For each row $i$ in $L_p$, variables of $A_{ii}$ that appear in $U_p$ are removed and we say that we perform **variable elimination in the row direction**. For each column $j$ in $U_p$, variables of $A_{ij}$ that appear in $L_p$ are removed. This will be referred to as **variable elimination in the column direction**. We then say that our algorithm performs **variable elimination in one direction**. Note that if, at a given step, variables are removed from both row $i$ and column $i$, it means that $i \in L_p$ and $i \in U_p$. In Section 4.3.1, we will prove that under additional assumptions more pruning of the variables could have been introduced. We then however comment in Section 4.3.2 that doing so makes reducibility detection as done in the “else” block of Algorithm 4.1 impossible. Note that this additional pruning would have improved the accuracy of our structural metrics as explained in Section 4.4.

The “else” block of Algorithm 4.1 shows that the algorithm can exploit the fact that when only one of the $L_p$ or $U_p$ is empty we can prune the current pivot from the quotient graph $G$ and therefore benefit from the reducibility of the input matrix. This feature of the QMLS algorithm will be justified in Section 4.3.2. We will also explain why it is correlated with the strategy used to prune variables.

**4.3.1. Two-way variable elimination**. Property 4.1 shows that under additional assumptions the structure of the quotient graph can be further pruned.

**Property 4.1.** If at step $k$, the entry at position $(i, j)$ is the only entry in row $i$ and column $j$ of $C^k$ then:

1. If $(i, j) \in L_p$, all the variables belonging to $L_p$ can then be removed from $A_{ij}$ (even though $j \notin U_p$),
(2) if $j \in U_p$, all the variables belonging to $U_p$ can then be removed from $A_p$, (even though $i \notin \mathcal{L}_p$).

Proof. From equation (2.2), we have $\text{Pattern}(C^{k+1}) \subseteq \text{Pattern}(C^k)$. Therefore, if at step $k$, $(i,j)$ is the only entry in row $i$ and column $j$ of $C^k$, it will remain the only entry in its row and column for all subsequent $C^l$, for $l > k$. Thus, it is for sure that $(i,j)$ will be selected as a pivot in a future step, and we can anticipate where local symmetrization will occur. So the entries in $A_{i,j} \cap \mathcal{L}_p$ for property 4.1(1) (or in $A_{i,j} \cap U_p$ for property 4.1(2)) can be pruned and will be retrieved from $\mathcal{L}_p$ (or $U_p$) when $(i,j)$ is eliminated. □

When we apply Property 4.1, we say that the algorithm performs elimination in both row and column directions. This process will be referred to as two-way variable elimination. For example when the pivot choice is limited to a transversal, the two-way variable elimination can be performed at each step of the elimination, as in the DMLS algorithm [4]. This is illustrated in Figure 4.1. We assume, for the sake of clarity, that the input matrix has been permuted to have all the candidate pivots on the diagonal. The shaded areas correspond to the variables that can be removed from the variable adjacency lists because they are implicitly stored through the adjacency lists of element $p$.

![Fig. 4.1. Illustration of two-way variable elimination.](image)

If the hypothesis of Property 4.1 is not true, the two-way variable elimination cannot be applied because we do not know whether local symmetrization will be performed or not. Consider Figure 4.1 again, if all three entries $(i,i)$, $(j,j)$ and $(j,i)$ belong in $C$, then we cannot prune all the shaded areas. This is because both $(i,i)$ and $(j,i)$ can be potential pivots from column $i$. If $(i,i)$ is chosen as the pivot from that column, then the shaded area in column $i$ can be pruned, because $i \in \mathcal{L}_p$ and local symmetrization is invoked. However, if $(j,i)$ is selected as the pivot from that column, then since $j \notin \mathcal{L}_p$ and $i \notin U_p$, the element $p$ will not be used to build the row and column adjacency of $(j,i)$, and the shaded area in column $i$ cannot be pruned. Otherwise, we cannot retrieve those variables from anywhere. The same observation applies to the shaded area in row $j$ if $(j,i)$ instead of $(j,j)$ is chosen as the pivot. Note that the shaded area in column $j$ can be pruned, because the entry $(p,j)$ is not a locally symmetrized entry, and so the variable elimination in the column direction can be applied.

4.3.2 Reducibility detection. If the input matrix is reducible, we may encounter a pivot $p$ such that either (1) both $\mathcal{L}_p = \emptyset$ and $U_p = \emptyset$ (we may call it
(strongly reducible) or (2) $\mathcal{L}_p = \emptyset$ or $\mathcal{U}_p = \emptyset$ (we may call it weakly reducible). Ideally, we would like to remove $p$ from the quotient graph $\mathcal{G}$ in both reducible cases. However, we will show that whether $p$ can be removed or not depends on whether we use only one-way variable elimination or use two-way variable elimination as well.

**Property 4.2.** If the variable elimination is always done in one direction, then the current pivot $p$ can be removed from the quotient graph if it is weakly reducible.

Property 4.2 comes from the fact that the pruning of the structures due to local symmetrization has not been anticipated. Thus, none of the entries in $\mathcal{L}_p$ (if $\mathcal{U}_p = \emptyset$) or $\mathcal{U}_p$ (if $\mathcal{L}_p = \emptyset$) will be needed by the other variables to represent their adjacency structure in $\mathcal{G}$. Therefore, pivot $p$ can be removed from $\mathcal{G}$.

**Property 4.3.** If the two-way variable elimination has been done at least once, then the current pivot $p$ can be safely removed from the quotient graph if and only if it is strongly reducible.

**Proof.** Firstly, when $\mathcal{U}_p = \emptyset$ and $\mathcal{L}_p = \emptyset$, $p$ becomes a singleton element and can certainly be removed from the quotient graph. Secondly, let us suppose that two-way variable elimination has been performed at least once. We now build a counter example to show that we cannot safely (in all possible cases) remove a weakly reducible pivot $p$. Let us assume without loss of generality that $\mathcal{U}_p = \emptyset$ and $\mathcal{L}_p \neq \emptyset$. Let us assume that there exists a variable $i \in \mathcal{L}_p$ and that there is an entry $(i,j)$ that is the only entry in row $i$ and column $j$ in $C$. We also assume that variables in $\mathcal{A}_{ij}^{(0)}$ have been pruned under two-way variable elimination. Therefore, $p$ must be used to retrieve those entries and cannot be removed from $\mathcal{G}$. This is illustrated in Figure 4.2 where the shaded area (1) in column $j$ is first stored through element $e$ then stored through element $p$ (after pivot $p$ absorbs element $e$). \[\square\]

Property 4.3 indicates a drawback of the two-way variable elimination: we can only prune the pivot in the strongly reducible case. The algorithm may be very inefficient if the matrix is very reducible in the weak sense.

![Fig. 4.2. Effect of variable elimination in both directions on reducibility detection (S indicates the position of local symmetrization).](image)

In the rest of this section, we discuss the detection and/or the exploitation of the BTF (Block Triangular Form) of the initial matrix. We show that, under certain assumptions, using one-way variable elimination, we can exploit and detect part of the BTF.

**Definition 4.1.** We say that an ordering is **BTF compatible** when

1. it selects only pivots within the diagonal blocks of the BTF;
   2. it eliminates all the pivots in a block before processing another block;

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(3) the order in which the diagonal blocks are considered must be such that each selected block is not adjacent to any remaining block in either the row or the column direction.

Note that condition (3) allows a block to be considered in the order of the BTF but also allows some slight variation of this form as will be shown in Figure 4.3.

Property 4.4. If CMVS selects pivots with a BTF compatible ordering then the sparsity of the block triangular form will be fully preserved and CMVS will generate a forest in which each tree corresponds to an irreducible component of the BTF.

This property is a consequence of Property 4.2. We illustrate in the following why condition (3) is needed in Definition 4.1. We also comment on the benefit of using one-way variable elimination when the matrices are reducible and on the adverse effect of two-way variable elimination for reducible matrices.

Let us assume that the constraint matrix $C$ is restricted to the diagonal of $A$ and that pivots are chosen down the diagonal in order. In this case, the two-way variable elimination could also be applied at each step of the ordering (see Property 4.1). Figure 4.3 shows a matrix with five full diagonal blocks of size $k$. The original matrix in BTF is shown on the left-hand side. The structure resulting from the use of a modified CMVS algorithm that would include two-way variables elimination is shown in the middle of the figure. We see that, with two-way variable elimination, local symmetrization interconnects all the diagonal blocks of the BTF and the BTF forest (independence between the diagonal blocks) is lost. In addition, because of the nonzero “s” between block (1) and block (2), the first row of block (2) is completely filled, which in turn will create fill-ins in all the lines of block (2) and in all the lines of the subsequent blocks.

With only one-way variable elimination, the last pivot of block (1) will be removed (see Property 4.2). The same holds for all subsequent blocks so that Algorithm 4.1 does not introduce any fill-in and produces a forest in which each tree corresponds to a diagonal block of the BTF matrix.

![Figure 4.3](image)

**Fig. 4.3. Illustration of two BTF compatible orderings of a matrix $A$ in BTF. Using two-way variable elimination introduces fill-ins in the factors.**

We now assume that the blocks of the matrix have been symmetrically permuted in the order $1,5,3,4$, see the right-hand side matrix in Figure 4.3. This ordering does not respect the BTF, however, it is BTF compatible thanks to the condition (3) of Definition 4.1. Property 4.2 can be applied to the last pivot of each block, since either $U_p$ or $L_p$ will be empty. The last pivot selected in each block will become the root of the tree associated with that block of the BTF matrix. Algorithm 4.1 thus fully exploits the BTF and preserves the sparsity of matrix. On the other hand, if we use two-way variable elimination, because of the “s” entry at position $(k+1,k)$, row
$k + 1$ becomes completely filled and after eliminating the $(k + 1)$st diagonal entry the remaining submatrix becomes also full.

![Diagram](image-url)

**Fig. 4.4.** The initial matrix is not in BTF (left plot) and the ordering is not BTF compatible (factors in the middle plot). The ordering on the right is BTF compatible and preserves sparsity. F: fill-in. S: local symmetrization.

Finally, we show in Figure 4.4 that when the ordering of the blocks does not satisfy condition (3) of Definition 4.1, CMS algorithm may partially lose BTF. In this example, local symmetrization between blocks (1) and (2) will create a dependency between them and fills one rectangular block because of the propagation of this dependency. Block (3) will however remain an independent tree because the last pivot $p$ in block (2) will have $\mathcal{U}_p = \emptyset$ and $p$ will become the root of the tree containing blocks (1) and (2). Note that if two-way variable elimination were applied, the ordering would not have detected the independence of block (3).

### 4.4. Update of the structural metric

In this section, we describe two classes of local heuristics to estimate the structural quality of a pivot. In the preamble section, we first describe the common framework that is independent of the metric used. In Section 4.4.2, the approximate Markowitz cost [4] is briefly introduced. In Section 4.4.3, we describe a metric based on an upper bound on the fill-ins introduced at each step of elimination. This approximation of the fill-ins has been studied by the authors of AMD [1] for symmetric matrices. We provide a generalization of this approximation to unsymmetric matrices and prove that it is a tighter upper bound on the fill-ins than the approximations proposed for symmetric matrices in [33]. Note that concerning the deficiency approximation in [30], there is no guarantee that it is an upper bound of the fill-in. Our approximate minimum fill-in heuristic will be referred to as AMF1. Finally, in Section 4.4.4, we discuss the complexity of the algorithms used to update the structural metrics.

#### 4.4.1. Preamble

Let us assume that the $p^{th}$ pivot $(\text{row}_p, \text{col}_p)$ has been selected. All the entries in $(\mathcal{L}_p \times \mathcal{V}_c \cup \mathcal{V}_r \times \mathcal{U}_p) \cap \text{Pattern}(C)$ are involved in the structural metric updates since the structure of rows in $\mathcal{L}_p$ and columns in $\mathcal{U}_p$ might have changed. The size of this area is thus larger than the area involved in the update of the structure of $C$ since columns (resp. rows) of $\mathcal{U}_p \setminus \text{R}_{\text{row}_p}$ (resp. of $\mathcal{L}_p \setminus \text{C}_{\text{col}_p}$) need also be considered. The algorithm to update the structural metrics will be one of the most costly steps of our algorithm.

We want the metrics to reflect the structural quality of an entry if it were selected as the next pivot. That is why we compute metrics which are related to the structure of our quotient graph and for which local symmetrization has been applied. In the following, the degrees, approximate degrees, fill-ins and approximate fill-ins are all related to this quotient graph structure. Note that, since for symmetric matrices our quotient graph becomes the standard symmetric quotient graph, all the discussions naturally apply to symmetric matrices.
Equations (3.1) and (3.2) that include local symmetrization could be used to compute the exact external degrees, but it would be costly. Instead, similar to AMD [1] and DMLS [4] algorithms, approximate row and columns external degrees can be computed. The AMD like approximate external row and columns degree, $amd_r(i, j)$ and $amd_c(i, j)$ respectively, are defined by the following two equations:

$$
amd_r(i, j) = |A_{i*} \setminus U_p| + |U_p \setminus J| + \sum_{e \in R_j \setminus C_j} (|U_e \setminus U_p|) + \sum_{e \in C_j \setminus R_j} (|U_e \setminus U_p|) - \alpha_j,
$$

with $\alpha_j = \max(|C_j|, 1)$ if $j \notin U_p$ else $\alpha_j = 0$.

(4.1)

$$
amd_c(i, j) = |A_{*j} \setminus L_p| + |L_p \setminus i| + \sum_{e \in R_i \setminus C_j} (|L_e \setminus L_p|) + \sum_{e \in C_j \setminus R_i} (|L_e \setminus L_p|) - \beta_i,
$$

with $\beta_i = \max(|R_i|, 1)$ if $i \notin L_p$ else $\beta_i = 0$.

(4.2)

As observed in [4], degree corrections ($\alpha_j$ and $\beta_i$ in equations (4.1) and (4.2)) are introduced to improve the approximations of the row and column external degrees in the presence of local symmetrization. To justify these correction terms, one can observe that if $j \notin U_p$ then $j$ is counted in every $U_e \setminus U_p$ for $e$ that is adjacent to column $j$ ($e \in C_j$). Furthermore, if $C_j$ is empty and $j \notin U_p$ then column $j$ has been counted in $A_{*j} \setminus U_p$ and should then be subtracted. This explains the use of $\alpha_j$ in the correction. $\beta_i$ can be justified in a similar way. The $|U_e \setminus U_p|$ and $|L_e \setminus L_p|$ quantities are computed similarly as in the AMD algorithm.

Note that since only one-way variable elimination is employed, the computation of the metric is less accurate than with two-way variable elimination. This is because in the latter case, for any element $e$, row index $i \in U_p$ and column index $j \in L_p$, we have $A_{i*} \cap U_e = \emptyset$ and $A_{*j} \cap L_e = \emptyset$. This is no longer true when one-way variable elimination is used (see Algorithm 4.1). But as was shown in Section 4.3, the benefit of one-way variable elimination is to better exploit the BTF of the matrix.

### 4.4.2. Approximation of Markowitz cost.

The approximation of the Markowitz cost comes directly from our approximation of the external row and column degrees. After eliminating the $k^{th}$ pivot, we define

$$
\overline{amd}_r(i, j) = \min(amd_r(i,j), n - k - 1),
$$

and

$$
\overline{amd}_c(i, j) = \min(amd_c(i,j), n - k - 1).
$$

The metric associated with the approximate Markowitz cost is then defined as

$$
\text{metric}^{(k+1)}(i, j) = \min \left\{ \overline{amd}_r(i, j) \times \overline{amd}_c(i, j), \left| \frac{1}{U_p \setminus j} \times \overline{amd}_c(i, j) \right|, \left| \frac{1}{L_p \setminus i} \times \overline{amd}_r(i, j) \right|, \left| \frac{1}{U_p \setminus j} \times \frac{1}{L_p \setminus i} \right| \right\}
$$

(4.5)

Here, we use the convention that if $(i, j)$ is a new entry in $C$, then $\text{metric}^{(k)}(i, j)$ in equation (4.5) is set to $+\infty$. Note that during the update of the degrees, contrary to AMD, we do not use the values of approximate row and column degrees computed at the previous step. This could have been done but would have required us to store two other arrays (one for row degrees, one for column degrees) of size $|C|$ since the approximate degree varies within a given row or column.
4.4.3. Approximation of the fill-in. With a minimum fill-in based metric, we want to estimate the new fill-in that would occur in the reduced matrix if an entry were selected as the next pivot. For the sake of completeness, one should mention that the fill-in metric of variables at distance two from the pivot might also vary. In this work, we will only consider variables adjacent to the pivot, i.e. at distance one, since results on symmetric matrices have shown that considering distance two variables significantly increases the complexity of the algorithm for relatively little gain in the quality of the ordering [30, 33].

A coarse upper bound of the fill-in that would occur can be obtained by removing the area corresponding to $\mathcal{L}_p \times \mathcal{U}_p$ from the Markowitz cost or the area corresponding to the largest adjacent clique [33]. A tighter approximation of the fill-in in the factors can be obtained by removing all the areas already filled during the elimination of the previous elements. We explain how to compute this new upper bound using already computed information and local correction terms.

Suppose that $i \in \mathcal{L}_p$ or $j \in \mathcal{U}_p$. Let $e$ be an element that belongs to $\mathcal{F}$. To simplify the notation we define $\bar{\mathcal{L}}_e = (\mathcal{L}_e \setminus \mathcal{L}_p) \setminus \{i\}$, $\bar{\mathcal{U}}_e = (\mathcal{U}_e \setminus \mathcal{U}_p) \setminus \{j\}$, $\bar{A}_e = (A_e \setminus \mathcal{L}_p) \setminus \{i\}$, $\bar{A}_u = (A_e \setminus \mathcal{U}_p) \setminus \{j\}$, $\bar{\mathcal{L}}_e = \mathcal{L}_e \setminus \{i\}$ and $\bar{\mathcal{U}}_e = \mathcal{U}_e \setminus \{j\}$. Thus a $\bar{}$ means that we subtract the current pivot structure and the current variable and a $\bar{}$ means that we only subtract the current variable.

Let $d_r(i, j)$ and $d_c(i, j)$ denote the external row and column external degrees of entry $(i, j)$ respectively. With our notation we have:

\begin{align}
\text{(4.6)} & \quad d_r(i, j) = |(\mathcal{U}_p \cup A_{i*} \cup \bigcup_{e \in \mathcal{F}} \mathcal{U}_e) \setminus \{j\}| \leq |\bar{\mathcal{U}}_p| + |\bar{A}_{i*}| + \bigcup_{e \in \mathcal{F}} |\bar{\mathcal{U}}_e| \\
\text{(4.7)} & \quad d_c(i, j) = |(\mathcal{L}_p \cup A_{j*} \cup \bigcup_{e \in \mathcal{F}} \mathcal{L}_e) \setminus \{i\}| \leq |\bar{\mathcal{L}}_p| + |\bar{A}_{j*}| + \bigcup_{e \in \mathcal{F}} |\bar{\mathcal{L}}_e|.
\end{align}

Let $S(i, j)$ denote the union of the areas associated with all the elements adjacent to entry $(i, j)$:

\[ S(i, j) = \bigcup_{e \in \mathcal{F}} (\bar{\mathcal{L}}_e \times \bar{\mathcal{U}}_e) \setminus (\mathcal{L}_p \times \mathcal{U}_p). \]

Ideally one might want to subtract both $|\bar{\mathcal{L}}_p \times \bar{\mathcal{U}}_p|$ and $S(i, j)$ from the Markowitz cost. An upper bound of the fill-in that would occur (including local symmetrization) if an entry $(i, j)$ were eliminated is:

\[ d_r(i, j)d_c(i, j) - |\bar{\mathcal{U}}_p||\bar{\mathcal{L}}_p| - S(i, j). \]

The authors of [33] have observed that instead of using the exact external degrees one could use the approximate (in the sense of the AMD algorithm) external degrees since both produce results of comparable quality and since AMD based metrics are significantly faster to compute. In this context, the corresponding upper bound of the fill-in metric becomes

\begin{align}
\text{(4.8)} & \quad \text{amd}_r(i, j)\text{amd}_c(i, j) - |\bar{\mathcal{U}}_p||\bar{\mathcal{L}}_p| - S(i, j).
\end{align}

Let $AS$ be an overestimation of area $S$,

\begin{align}
\text{(4.9)} & \quad AS(i, j) = \bigcup_{e \in \mathcal{F}} |(\bar{\mathcal{L}}_e \times \bar{\mathcal{U}}_e) \setminus (\mathcal{L}_p \times \mathcal{U}_p)| = \bigcup_{e \in \mathcal{F}} |\bar{\mathcal{U}}_e||\bar{\mathcal{L}}_e| + |\bar{\mathcal{L}}_e||\bar{\mathcal{U}}_e \cap \mathcal{U}_p|.
\end{align}
Property 4.5 proves that one can in fact subtract area $AS(i, j)$, instead of $S(i, j)$, from equation (4.8), to obtain a more accurate upper bound of the fill metric.

**Property 4.5.** $Amd_r(i, j)Amd_c(i, j) - |Up|\mid L_p\mid - AS(i, j)$ is an upper bound of the fill-in that would occur in the quotient graph if $(i, j)$ were eliminated.

An intuitive proof of Property 4.5 is that, during the computation of the approximate degree, the submatrix is expanded in such a way that the intersections between all $U_e$ and between all $L_e$ for $e \in \mathcal{F}$ are empty. The area $AS$ corresponds to a real surface in the expanded matrix and can be removed from the area $amdr(i, j)amdc(i, j)$ to compute the fill-in that would occur in the expanded matrix (see Figure 4.5). Moreover, this fill-in in the expanded matrix is an upper bound of the exact fill-in in the quotient graph. A formal proof of Property 4.5 is given in [32].

We now explain that, although computing $AS(i, j)$ is not trivial it is however not costly. To compute the area $AS$, we need to evaluate $U_e$, $L_e$, $U_p$ and $L_p$ (see equation (4.9)). It appears to be difficult to compute these quantities directly. Indeed only the quantities $|U_e|$, $|L_e|$, $|U_e \setminus U_p|$, $|L_e \setminus L_p|$, $|U_e \cap U_p|$ and $|L_e \cap L_p|$ are known thanks to $AMD$ like computations. For each entry $(i, j)$ involved in the metric update, two quantities are computed: an approximation of $AS(i, j)$, $area_{ij}$, such that

$$
area_{ij} = \sum_{e \in \mathcal{F}} (|U_e \setminus U_p|\mid L_e\mid + |L_e \setminus L_p|\mid U_e \cap U_p|),
$$

and local correction terms, $cor_{loc_{ij}}$, so that

$$
AS(i, j) = area_{ij} - cor_{loc_{ij}}.
$$

Computing $area_{ij}$ only involves already known quantities. To compute the local correction terms $cor_{loc_{ij}}$ we have to take into account for each $e \in \mathcal{F}$ all possible cases: $e \in R_i \setminus L_j$, $e \in L_j \setminus R_i$, and $e \in L_j \cap R_i$. 

---

**Figure 4.5.** AMD1 areas in the expanded matrix.

| Fill of e1 (fill of p removed) | Fill of e2 (fill of p removed) |
Algorithm 4.2 explains how to compute both \( area_{ij} \) and \( corr_{ij} \) from which the AMFI metric can be deduced. To help understand how the local correction terms are computed, we indicate in Figure 4.6 the areas corresponding to different values of the local correction. In the following, we focus on the local correction terms computed during Loop 1 of the algorithm. For \( i \in \mathcal{L}_p \) (Loop 1), we note that, for each element \( e \in \mathcal{R}_i \), \( [\mathcal{U}_e \setminus \mathcal{U}_p] \) must be added to the correction term since it was taken into account by mistake in \( [\mathcal{U}_e \setminus \mathcal{U}_p] \mathcal{L}_e \). In Figure 4.6, this situation occurs to cases i1, i2, i3 and i4 but not to cases i5 and i6 for which \( [\mathcal{U}_e \setminus \mathcal{U}_p] \) is not added at line (1) of the algorithm since \( e \notin \mathcal{R}_e \) and \( e \notin \mathcal{R}_q \). We then see at line (2) of the algorithm that, since \( j \in \mathcal{U}_p \), \( [\mathcal{L}_e \setminus \mathcal{L}_p] \) must be added to the local correction (case i2). For case i3 at line [3] of the algorithm, the complete column structure \( [\mathcal{L}_e] \) has been taken into account in \( area_{ij} \) and must be added to the correction term. Furthermore, since the entry \( (i,j) \) had already been counted in \( [\mathcal{U}_e \setminus \mathcal{U}_p] \), one should remove it from the correction term at line [3] of the algorithm. For case i5 (i6) in Figure 4.6, we add \( [\mathcal{L}_e \setminus \mathcal{L}_p] \mathcal{L}_p \) \( (\mathcal{L}_e) \) that were counted by mistake in term \( [\mathcal{L}_e \setminus \mathcal{L}_p] \mathcal{U}_e \cap \mathcal{U}_p \) \( (\mathcal{L}_e \setminus \mathcal{U}_p] \mathcal{L}_e \) of \( area_{ij} \).

In practice we use \( amd_r(i,j) \) and \( amd_c(i,j) \) as defined in equations (4.3) and (4.4) instead of \( amd_r(i,j) \) and \( amd_c(i,j) \). Because of that it may happen that \( amd_r(i,j) amd_c(i,j) - [\mathcal{U}_p \setminus j] \mathcal{L}_p - AS(i,j) \) becomes negative, meaning that either \( amd_r(i,j) < amd_r(i,j) or amd_c(i,j) < amd_c(i,j) \). In such cases, as is done in AMD and DMLS, one can artificially set the metric to 0. We propose here an alternative, that could also be applied to these approaches to limit the tie-breaking. We introduce row and column scaling terms

\[
\text{rowscale} = \frac{amr_r(i,j)}{amd_r(i,j)} \quad \text{and} \quad \text{colscale} = \frac{amr_c(i,j)}{amd_c(i,j)}.
\]

If one systematically scales the area \( AS \) by \( \text{rowscale} \times \text{colscale} \), then we ensure a positive metric and avoid tie-breaking problems due to metrics equal to 0. Our final
Algorithm 4.2 Computation of areas and local corrections for AMFI metric.

for $i \in L_p$ do /* Loop 1 */
  area$\_savei = \sum_{e \in R_i} (|\mathcal{L}_e \setminus \mathcal{U}_p| + |\mathcal{L}_p \setminus \mathcal{U}_e|)$
  cor$\_loc$save$\_i = \sum_{e \in R_i} |\mathcal{L}_e \setminus \mathcal{U}_p| /* cases i1, i2, i3 and i4 */
for each $j$ such that $(i, j) \in C$ do
  area$_{ij} = area$\_save$\_i$, cor$\_loc$$_{ij} = cor$\_loc$save$\_i$
for each $e \in C_j$ do
  if $e \in R_i$ then /* $e$ already visited, its fill-in already counted in area$\_ij */
    if $j \in U_p$ then
      cor$\_loc$$_{ij} = cor$\_loc$$_{ij} + |\mathcal{L}_e \setminus \mathcal{U}_p| /* case i2 */
    else /* cases i3 and i4 */
      cor$\_loc$$_{ij} = cor$\_loc$$_{ij} + |\mathcal{L}_e \setminus \mathcal{U}_p| - 1 /* case i3 */
  end if
  else /* $e$ not already visited, we need to count its corresponding area */
    area$\_ij = area$\_save$\_i$ + $|\mathcal{L}_e \setminus \mathcal{U}_p| + |\mathcal{L}_e \setminus \mathcal{L}_p| + |\mathcal{L}_p \setminus \mathcal{U}_e|$
    if $j \in U_p$ then
      cor$\_loc$$_{ij} = cor$\_loc$$_{ij} + |\mathcal{L}_e \setminus \mathcal{L}_p| /* case i5 */
    else /* cases i6 and i7 */
      cor$\_loc$$_{ij} = cor$\_loc$$_{ij} + |\mathcal{L}_e | /* case i6 */
      end if
  end if
end for
for each $j \in U_p$ do /* Loop 2 */
  area$\_savej = \sum_{e \in C_j} (|\mathcal{L}_e \setminus \mathcal{U}_p| + |\mathcal{L}_p \setminus \mathcal{U}_e|)
  cor$\_loc$save$\_j = \sum_{e \in C_j} |\mathcal{L}_e \setminus \mathcal{U}_p| /* used for cases 8 and 9 */
for each row $i$ such that $(i, j) \in C$ and $i \notin L_p$ do
  area$\_ij = area$\_save$\_j$, cor$\_loc$$_{ij} = cor$\_loc$save$\_j$
  for each $e \in R_i$ do
    if $e \notin C_j$ then /* cases i8 and i9 */
      area$\_ij = area$\_save$\_j$, cor$\_loc$$_{ij} = cor$\_loc$save$\_j$
      area$\_ij = area$\_save$\_j$, cor$\_loc$$_{ij} = cor$\_loc$save$\_j$
      end if
    else /* cases i8 and i9 */
      area$\_ij = area$\_save$\_j$, cor$\_loc$$_{ij} = cor$\_loc$save$\_j$
      end if
  end for
end for
end for

AMFI metric is then defined as follows:

\[
\text{metric}^{(k+1)}(i, j) = \min \left\{ \begin{array}{l}
\text{amdi}(i, j) - |\mathcal{U}_p \setminus j| |\mathcal{L}_p \setminus i| \\
- \text{rrouscale} \times \text{colscale} \times \text{AS}(i, j)
\end{array} \right\}
\]

\[
(4.12) \quad \text{metric}^{(k+1)}(i, j) = \min \left\{ \begin{array}{l}
\text{amdi}(i, j) - |\mathcal{U}_p \setminus j| |\mathcal{L}_p \setminus i| \\
- \text{rrouscale} \times \text{colscale} \times \text{AS}(i, j)
\end{array} \right\}
\]

4.4.4. Complexity. First note that the complexity of the update of all the proposed metrics is comparable since the dominant cost associated with the computation of the approximate row and column degrees is shared by all metric calculations. We will thus focus here on the computation of the approximate Markowitz cost and more particularly on the computation of $\text{amdi}_c$ as defined by equation (4.2). The quantity $|A_{ij} \setminus \mathcal{L}_p|$ is computed only once for all the entries in column $j$ belonging to $C$. The computation time is thus at each step bounded by $O(\sum_j |A_{ij}|) = O(|E|)$. For each row in $L_p$, the computational cost of the third term of Equation (4.2) is $O(|R_i|)$. For each row in $L_p$ and each column $j$ such that $(i, j)$
belongs to $\mathbf{C}$, the computational cost of the fourth term of equation (4.2) is $O(|C_j|)$. Thus the total complexity of $amdc$ for all the $(i,j)$ entries in $(\mathcal{L}_p \times \mathcal{V}_p) \cap \mathcal{P}attern(\mathbf{C})$ is

\begin{equation}
O \left( \sum_{i \in \mathcal{L}_p} |R_i| + \sum_{i \in \mathcal{L}_p} \sum_{j \in R_i} |C_j| \right).
\end{equation}

Let us define $n_1, \ldots, n_{|\mathcal{L}_p|}$ by induction: $n_1 = \arg \max |C_j|$ and for $i > 1$, $n_i = \arg \max_{j \neq n_{i-1}, \ldots, n_1} |C_j|$. At each step $p$ of the symbolic factorization, Let us define $k^p$ as the largest row or column length. $k^p = \max_{i \in \mathcal{V}_p} |R^i_p|, \max_{j \in \mathcal{V}_p} |C^j_p|$). The double sum of equation (4.13) thus contains at most $k^p|\mathcal{L}_p|$ terms where $k^p$ is defined as the size of the largest row/column at step $p$. Note that at most $k^p$ terms can be equal to $C_{n_i}$ for all $i \in [1, |\mathcal{L}_p|]$. Thus, we have the following inequality: $\sum_{i \in \mathcal{L}_p} \sum_{j \in R_i} |C_j| \leq k^p \times \sum_{1 \leq i \leq |\mathcal{L}_p|} |C_{n_i}|$. Then we have

$$\sum_{i \in \mathcal{L}_p} (|R_i| + \sum_{j \in R_i} |C_{n_j}|) \leq k^p \times \left( \sum_{i \in \mathcal{L}_p} |R_i| + \sum_{1 \leq i \leq |\mathcal{L}_p|} |C_{n_i}| \right).$$

Finally, thanks to the in-place property of our algorithm to update the bipartite quotient graph we obtain

$$\sum_{i \in \mathcal{L}_p} |R_i| + \sum_{1 \leq i \leq |\mathcal{L}_p|} |C_{n_i}| \leq |E|.$$

The same bounds can be obtained for $amdc$, so that the total cost for updating the metric is $O(k^p|E|)$ per elimination step. The total time complexity is $O(\sum_{1 \leq p \leq n} k^p|E|)$. If $k^p$ is bounded, the complexity becomes $O(n|E|)$ which is the same complexity as the AMD algorithm [26].

4.5. **Supervariables and mass elimination.** For the sake of clarity, the algorithms described in the previous section did not include supervariables. In this section, we first define our generalization of supervariables and mass elimination to bipartite quotient graphs with off-diagonal pivots. We then revisit the previous algorithms and explain what has to be modified to detect and exploit supervariables.

4.5.1. **Adaptation of CMLS main scheme.** In our context, we want supervariables to exploit identical adjacency structures in the graph at each step of the elimination. Supervariables are thus defined on the bipartite quotient graph of $\mathbf{A}$, whereas on the bipartite graph of $\mathbf{C}$ we only use simple variables. With the CMLS algorithm we cannot use exactly the same kinds of supervariables as in [1, 4, 16, 21] because they assume that pivots are on the diagonal so that a row can be associated with a column before being selected as pivot. That is why our concept of supervariable is closer to the one used in [20]: we define *indistinguishable row variables* (resp. *indistinguishable column variables*) as row variables (resp. columns variables) which have the same adjacency in $\mathcal{G}$. To limit the cost of supervariable detection, two hash functions (see for example [7]) are then used for each row and column direction.

If $i$ and $j$ are two indistinguishable row variables, they are replaced in $\mathcal{G}$ by a row *supervariable* containing both $i$ and $j$, labeled by its *principal row variable* ($i$, say) [15, 16, 17]. The notation $i$ is used to denote this row supervariable and $i = \{i, j\}$. $i$ and $j$ are said to be *constituent row variables* of the row supervariable $i$ and the notations $i \in$
i and j ∈ i are then used. At the beginning of Gaussian elimination, the row variables are said to be \textit{simple row variables}. Each simple row variable i can also be seen as a row supervariable i = \{i\}. For each row supervariable i, |i| corresponds to its size, i.e., its number of constituent variables. Similar definitions and notation can be introduced for the column supervariables, the \textit{principal column variables}, the \textit{constituent column variables} and the \textit{simple column variables}. When it is clear from the context, we do not differentiate between a column or a row supervariable. Furthermore, let r1 and r2 be two row variables which belong to the same row supervariable r and c1 and c2 be two column variables which belong to the same column supervariable c. After the pivot \( p_1 = (r_1, c_1) \) is eliminated, \( p_2 = (r_2, c_2) \) can be eliminated in \( G \) without causing extra fill-in. This process, commonly referred to as \textit{mass elimination} [23], creates a new (super)element \( e = (r, c) \) in the quotient graph. In the following, we comment on the algorithmic modifications due to the introduction of supervariables.

Let \( p = (r, c) \) be the current pivot. The first modification of the algorithm concerns the introduction of a scaling of the structural metrics as defined either by equation (4.5) or by equation (4.12). The structural metric of an entry \((i,j)\) adjacent to \( p \) either in the row or column direction is divided by \( \min(|i|,|j|) \). Indeed \( \min(|i|,|j|) \) corresponds to the size of the largest pivot block which could be eliminated if a pivot at the intersection of these row and column supervariables were selected.

The second modification of the algorithm concerns the elimination process which is performed in the three main steps. During the first step, the scaled metric is use to select a pivot in \( C \). During the second step, we retrieve its associated row \( r \) and column \( c \) supervariable in \( A \). During the third step, we eliminate “as many as possible” variables belonging to \( (r \times c) \cap C \). Note that the meaning of “as many as possible” will depend on the context. If a hybrid strategy is used then pivot entries might be rejected because of numerical criteria. Furthermore, since the \( C \) matrix is updated when eliminating a pivot, the new nonzero entries that might be in the intersection of the pattern of \( C \) and the supervariables need also be considered. The same modified three steps are also applied to the mass elimination process. Finally, if some constituent variables of a supervariable have not been eliminated, then they are built to build a new supervariable and are re-inserted in \( G \).

The final modification concerns the update of the structural metric that will be fully described in Section 4.5.2. Note that during the metric update, we also mark candidates for mass elimination.

\textbf{4.5.2. Revisiting computation of the structural metrics.} When using supervariables and after the elimination of a pivot \( p \), the approximate external row and column degrees as defined by equations (4.1) and (4.2) become:

\begin{equation}
\text{amd}_i(i,j) = |A_i \setminus U_p| + |U_p \setminus \{j\}| + \sum_{e \in R_i \cup L_p} (|U_e \setminus U_p|) - \alpha_j |j|, \text{ with } \alpha_j = \max(|C_j|, 1) \text{ if } j \notin U_p \text{ else } \alpha_j = 0.
\end{equation}

\begin{equation}
\text{amd}_c(i,j) = |A_j \setminus L_p| + |L_p \setminus \{i\}| + \sum_{e \in R_j \cup L_p} (|L_e \setminus L_p|) - \beta_i |i|, \text{ with } \beta_j = \max(|R_i|, 1) \text{ if } i \notin L_p \text{ else } \beta_j = 0.
\end{equation}

In Section 4.4.3, we have shown that to evaluate the AMFI metric of an entry \((i,j)\) we need to compute the area \( AS(i,j) \) of equation (4.12) as the sum of two terms \( area_{ij} \) and \( cor_{ij} \) (see equation (4.11)). The \( area_{ij} \) term (see equation (4.10)) is related to the areas of all the bi-cliques of the already eliminated elements so that its computation does not depend on the use of supervariables. The computation of the
local correction terms $cor_{loc_{ij}}$ must however be revisited. We split this correction term into two parts. Each part, $row_{cor_{loc_{ij}}}$ and $col_{cor_{loc_{ij}}}$, refers respectively to entries in the row supervariable and to entries in the column supervariable so that the new equation for defining the area $AS$ is:

$$AS(i,j) = area_{ij} - row_{cor_{loc_{ij}}} |i| - col_{cor_{loc_{ij}}} |j|.$$  

For an entry $(i,j) \in C$, the local corrections for the different cases of Figure 4.6 are then:

- Case 1: $row_{cor_{loc_{ij}}} = |U_c \setminus U_p|$ and $col_{cor_{loc_{ij}}} = 0$,
- Case 2: $row_{cor_{loc_{ij}}} = |U_c \setminus U_p|$ and $col_{cor_{loc_{ij}}} = |L_c \setminus L_p|$,
- Case 3: $row_{cor_{loc_{ij}}} = |U_c \setminus U_p|$ and $col_{cor_{loc_{ij}}} = |L_c| - |i|$,
- Case 4: $row_{cor_{loc_{ij}}} = |L_c \setminus L_p|$ and $col_{cor_{loc_{ij}}} = 0$,
- Case 5: $row_{cor_{loc_{ij}}} = 0$ and $col_{cor_{loc_{ij}}} = |L_c \setminus L_p|$,
- Case 6: $row_{cor_{loc_{ij}}} = 0$ and $col_{cor_{loc_{ij}}} = |L_c|$,
- Case 7: $row_{cor_{loc_{ij}}} = |U_c|$ and $col_{cor_{loc_{ij}}} = 0$,
- Case 8: $row_{cor_{loc_{ij}}} = |U_c| - |j|$ and $col_{cor_{loc_{ij}}} = |L_c \setminus L_p|$,
- Case 9: $row_{cor_{loc_{ij}}} = 0$ and $col_{cor_{loc_{ij}}} = |L_c \setminus L_p|$.

It is then straightforward to accumulate these corrections by adapting Algorithm 4.2.

5. Experiments. In this section we analyze the effect of the CMLS ordering on the performance of sparse solvers. Our new ordering will be compared to the combination of DMLS ordering and MC64 [13, 14] because it is the most robust in-place local heuristic (better than AMD, see [4]) in terms of numerical stability and fill-in reduction in the factors. DMLS takes into account the asymmetry of the matrices, selects pivots on the diagonal, applies local symmetrization and two-way variable elimination. Thus it can be considered a restricted CMLS. We recall that MC64 permutes the matrix such that the product of the diagonal elements is maximized.

With the CMLS ordering, our pivot sequence results from a combination of structural and numerical information (even when only structural metrics are used to select the pivots, the initialization of our constraint matrix is based on numerical considerations). Therefore it is important to analyze the numerical quality of the proposed sequence of pivots. In this context, for very different motivations, we may want to experiment with both an approach that performs partial pivoting to preserve numerical stability and an approach based on static pivoting. In the first case, the numerical quality of the proposed sequence of pivots is not so critical to obtaining a backward stable factorization and we expect to improve the sparsity of the factors because of the freedom to select entries in the constraint matrix $C$. In the case of a static pivoting, we expect that the capacity of CMLS to select pivots according to numerical criteria can be used to better control the numerical quality of the sequence while still offering some more freedom than a diagonal Markowitz algorithm. In fact with the CMLS algorithm we can define a family of orderings and expect that two probably different members of this family can be used in these two cases: a CMLS ordering in which $C$ offers a lot of freedom to choose the pivots and a CMLS ordering in which the selection of the pivots is strongly guided by the numerical values in $C$.

To represent each class of solver techniques, we consider the multifrontal code MA41_JENS [2, 5] which performs numerical pivoting during the factorization and the supernodal code SuperLUDIST [28] which performs static pivoting. Both codes are run in sequential mode. As shown in [3, 5, 10, 25] the approaches used to factorize the matrix in MA41_JENS and SuperLUDIST are very competitive in shared/sequential and distributed memory environments respectively. Note that because of the important
algorithmic similarities between \texttt{NA41JNS} and the distributed memory code \texttt{MUMPS},
this work will be also very beneficial to the distributed memory multifrontal code.

In Section 5.1, we present our experimental environment. In Section 5.2 we first
briefly discuss the case where the pivot choice in \texttt{OMLS} is restricted to the matching
provided by \texttt{MC64}. In Section 5.3, we analyze the behavior of our ordering when
a structural strategy is used to select the pivots. We report performance obtained
with \texttt{NA41JNS} solver in terms of time and memory used during factorization. In
Section 5.4, we illustrate the benefits resulting from the use of hybrid strategies for
pivot selection in the SuperLU_DIST context and focus in the case on the numerical
effects.

5.1. Experimental environment.

5.1.1. Test matrices and computing environment. A representative set
of 19 large unsymmetric matrices has been selected (see Table 5.1) from Tim
Davis' collection [9]. All our results have been obtained on a Linux PC computer
(Pentium 4, 2.8 GHz, 2 GBytes of memory and 1 MByte of cache). We use the
Portland Fortran 90 compiler pgf90, C compiler gcc (both with -O3 option) and
GOTO BLAS [24].

Consider a matrix $A = (a_{ij})$ and let $nnz(A)$ be its number of nonzero entries.
We define the structural symmetry $s(A)$ as:

\[
s(A) = \frac{\text{size} \{(i,j) \text{ s.t. } a_{ij} \neq 0 \text{ and } a_{ji} \neq 0\}}{\text{nnz}(A)}.
\]

In the remainder of this section, the symmetry of a matrix always refers to the
structural symmetry once the \texttt{MC64} permutation has been applied (see column $sym$
of Table 5.1).

<table>
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<tr>
<th>Group/Matrix</th>
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<th>$\text{nnz}$</th>
<th>$\text{sym}$</th>
<th>description</th>
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<td>1683942</td>
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<td>Single-material crack problem (sinco-bas0s)</td>
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</table>

We systematically apply random row and column permutations to our initial
matrix so that the ordering algorithms are less sensitive to the effects of tie-breaking,
We ran each problem with eleven random permutations and selected the run whose
ordering returns the median fill-in in the factors.

\[\text{Table 5.1}\]

Test matrices.
5.1.2. CMLS default parameters. The initialization of $C$ is done using scaled matrix and the maximum weighted matching returned by $\text{MP64}$. To limit the size of $C$ and the complexity (cost and memory) of the ordering phase, the initial number of entries in $C^0$ is set between $n$ and $4n$. We then drop the entries that are smaller than 0.1 in magnitude and the entries that have too large structural metrics. While dropping, we still maintain the property $M \subseteq \text{Pattern}(C)$. On our test set, we observed that the size of $C^0$ is between $n$ and $3n$ after this last dropping phase.

We use the Improved Approximate Minimum Fill metric $\text{AMFI}$ of Section 4.4.3 since it is the most efficient metric for both CMLS and DMLS orderings.

5.2. Preliminary remarks about diagonal constraint matrices. When $C_0$ contains only the entries from the $\text{MP64}$ matching and thus the set of candidate pivots for CMLS and DMLS is identical, one should expect a comparable behavior of the two algorithms in terms of fill-in in the factors. However, we have noticed (see [32] for detailed results) that CMLS ordering tends to produce sparser factors even if DMLS uses two-way variable elimination which leads (as explained in Section 4.4.1) to more accurate structural metrics. This can be explained by the following algorithmic differences:

- Thanks to the one-way variable elimination, CMLS can eliminate all the elements in both the strongly reducible and the weakly reducible situations. It has more chance to better preserve the sparsity of reducible matrices. This is well illustrated by the mult.decp matrix, which has 7448 irreducible components. DMLS and CMLS detect 875 singletons during a common preprocessing step. Then during ordering, DMLS detects 95 additional blocks versus 229 blocks for CMLS. Note that on the other matrices the number of detected blocks does not depend on the ordering and thus the algorithm used for variable elimination will not influence much the results on our set of test matrices.

- CMLS can create a row (column) supervariable if two rows (columns) have the same structure. DMLS can create a supervariable only if both rows and columns have the same structure. Thus, on the same quotient graph CMLS will detect more supervariables than DMLS. Note that the use of supervariables improves the accuracy of the structural metric. For example, if we consider that variables $i$ and $j$ belong to the same row supervariable, then the entries in $A_i*$ and $A_j*$ will not be counted as fill-in.

- In the CMLS implementation, we use rowscale and colscale coefficients (see end of Section 4.4.3) to reduce the amount of tie-breaking between variables that would have a negative metric (reset to 0) with DMLS.

We should stress that most of these algorithmic differences were justified because CMLS is designed to handle more general and complex situations than DMLS. What was not at all predicted is that the best implementation of DMLS should use the more general framework of the CMLS ordering.

5.3. Structural strategy.

5.3.1. Structure of the factors. In this section, we analyze the effect of the ordering on the size of the factors and compare the predicted size and the actual size of the factors. When there are no off-diagonal pivoting and node amalgamation, the actual size would be the same as the predicted size.

Table 5.2 compares CMLS with DMLS for both the estimated and the real size of the factors, using the $\text{MA41JNS}$ solver. For most matrices, the CMLS ordering results
in sparser factors. The gains in sparsity vary from 0% to 43%, with gains very much comparable for both the estimated and the real size of the factors.

<table>
<thead>
<tr>
<th>Matrix</th>
<th>Estimated size of factors</th>
<th>Real size of factors</th>
<th>Ratio (actual/predicted)</th>
</tr>
</thead>
<tbody>
<tr>
<td>av341002</td>
<td>6610</td>
<td>9359</td>
<td>0.69</td>
</tr>
<tr>
<td>g7jac300sc</td>
<td>28217</td>
<td>36345</td>
<td>29035</td>
</tr>
<tr>
<td>g7jac100sc</td>
<td>21561</td>
<td>31797</td>
<td>22222</td>
</tr>
<tr>
<td>ja99j9ac100sc</td>
<td>3793</td>
<td>3774</td>
<td>4190</td>
</tr>
<tr>
<td>ja99j9ac100sc</td>
<td>3712</td>
<td>3119</td>
<td>3505</td>
</tr>
<tr>
<td>lhr34c</td>
<td>6612</td>
<td>6833</td>
<td>7522</td>
</tr>
<tr>
<td>lhr734c</td>
<td>13958</td>
<td>13274</td>
<td>15860</td>
</tr>
<tr>
<td>mark3jac120ac</td>
<td>13079</td>
<td>13433</td>
<td>13668</td>
</tr>
<tr>
<td>mark3jac140ac</td>
<td>15127</td>
<td>15621</td>
<td>15790</td>
</tr>
<tr>
<td>bayer01</td>
<td>1374</td>
<td>2410</td>
<td>1533</td>
</tr>
<tr>
<td>sinc18</td>
<td>31314</td>
<td>37971</td>
<td>32382</td>
</tr>
<tr>
<td>sinc15</td>
<td>151453</td>
<td>16974</td>
<td>16093</td>
</tr>
<tr>
<td>Zhao2</td>
<td>12819</td>
<td>13037</td>
<td>13453</td>
</tr>
<tr>
<td>mult_dfop_03</td>
<td>780</td>
<td>865</td>
<td>959</td>
</tr>
<tr>
<td>twotone</td>
<td>7469</td>
<td>8152</td>
<td>8377</td>
</tr>
<tr>
<td>onetone1</td>
<td>2827</td>
<td>3174</td>
<td>3125</td>
</tr>
<tr>
<td>torso1</td>
<td>30505</td>
<td>30639</td>
<td>31491</td>
</tr>
<tr>
<td>bwmw1</td>
<td>37942</td>
<td>43382</td>
<td>35449</td>
</tr>
<tr>
<td>shermanACb</td>
<td>365</td>
<td>396</td>
<td>465</td>
</tr>
</tbody>
</table>

Mean CMLS/DMLS: 0.87 / 0.88  

| Table 5.2 | MA41JNS size of the factors and analysis reliability. Each number for the factor size is in thousands. |

As expected the fact that more flexibility has been offered to select off-diagonal pivots in the constraint matrix help CMLS at preserving the sparsity of the factors. Doing so we have however authorized CMLS to select pivots that do not belong to the maximum weighted matching. Since, a structural metric has then been used by CMLS algorithm to select the pivots, it is thus critical to evaluate the numerical quality of our sequence of pivots with MA41JNS. We recall that, thanks to partial threshold pivoting, the factorization phase of MA41JNS (the default value of the threshold is used in all experiments) will modify the pivot sequence to control the growth factor. This may result in an increase in the estimated factor size and number of operations. We thus also provide in Table 5.2 the ratio between the number of nonzeros in the factors and the forecast number of nonzeros in the factors (both delayed pivots and amalgamation contributes to increasing this ratio). Note that from a software point of view it is also critical for the estimation to reflect the reality. We see in Table 5.2 that the growth of the size of the factors is reasonable. Note that the fill-in due to the amalgamation process can be predicted and so if node amalgamation were included in the estimation, the growth of the size of the factors would be even smaller. Furthermore, it is even more interesting to observe that the growth obtained with CMLS is comparable to that obtained with DMLS. This shows that limiting the pivot choice to the maximum weighted matching (as done in DMLS) and considering numerical information in the constraint matrix (hybrid strategies to select pivots) are not critical in the context of a partial pivoting code.

5.3.2. Run-time and memory usage. In this section, we are concerned with the number of operations, run-time and memory usage of MA41JNS. The extra cost due to numerical pivoting during factorization is always included in the number of operations.
<table>
<thead>
<tr>
<th>Matrix</th>
<th>Number of operations</th>
<th>Memory needed</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CMS</td>
<td>DMIS</td>
</tr>
<tr>
<td>av4.1092</td>
<td>187.67</td>
<td>3809.2</td>
</tr>
<tr>
<td>g7jac3000c</td>
<td>27330.2</td>
<td>44379.8</td>
</tr>
<tr>
<td>g7jac1800c</td>
<td>17005.9</td>
<td>38434.8</td>
</tr>
<tr>
<td>jan99jac1200c</td>
<td>11623.5</td>
<td>1141.8</td>
</tr>
<tr>
<td>jan99jac1000c</td>
<td>1039.0</td>
<td>938.3</td>
</tr>
<tr>
<td>lhr71c</td>
<td>1099.4</td>
<td>3812.7</td>
</tr>
<tr>
<td>mark3jac1200c</td>
<td>3567.8</td>
<td>4981.5</td>
</tr>
<tr>
<td>mark3jac1400c</td>
<td>6786.6</td>
<td>7142.8</td>
</tr>
<tr>
<td>bayer01</td>
<td>58.9</td>
<td>155.4</td>
</tr>
<tr>
<td>sinc18</td>
<td>55501.4</td>
<td>77367.7</td>
</tr>
<tr>
<td>sinc15</td>
<td>19521.4</td>
<td>22388.5</td>
</tr>
<tr>
<td>Zhao2</td>
<td>8023.3</td>
<td>9250.6</td>
</tr>
<tr>
<td>multi_dcop.03</td>
<td>85.7</td>
<td>114.6</td>
</tr>
<tr>
<td>twotone</td>
<td>5142.9</td>
<td>5591.1</td>
</tr>
<tr>
<td>onetone1</td>
<td>1023.2</td>
<td>1250.3</td>
</tr>
<tr>
<td>tomo1</td>
<td>24141.6</td>
<td>25068.9</td>
</tr>
<tr>
<td>bbbmat</td>
<td>30455.0</td>
<td>40766.0</td>
</tr>
<tr>
<td>shermanACb</td>
<td>22.3</td>
<td>27.7</td>
</tr>
</tbody>
</table>

Table 5.3
MA14S number of operations (in millions) and memory used (in thousands of reals).

We see in Table 5.3 that on almost all the matrices, the CMS ordering reduces the amount of memory used, with an average reduction around 12%. The reduction in the number of operation is even more important (mean value of 21%) and will contribute to the reduction in the factorization time.

<table>
<thead>
<tr>
<th>Matrix</th>
<th>ordering time</th>
<th>factorization time</th>
<th>solution time</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CMS</td>
<td>DMIS</td>
<td>CMS</td>
</tr>
<tr>
<td>av4.1092</td>
<td>3.79</td>
<td>3.62</td>
<td>1.57</td>
</tr>
<tr>
<td>g7jac3000c</td>
<td>22.25</td>
<td>9.46</td>
<td>16.17</td>
</tr>
<tr>
<td>g7jac1800c</td>
<td>16.34</td>
<td>7.42</td>
<td>12.14</td>
</tr>
<tr>
<td>jan99jac1500c</td>
<td>6.07</td>
<td>2.59</td>
<td>1.68</td>
</tr>
<tr>
<td>jan99jac1000c</td>
<td>4.78</td>
<td>1.92</td>
<td>1.616</td>
</tr>
<tr>
<td>lhr71c</td>
<td>3.56</td>
<td>3.22</td>
<td>11.85</td>
</tr>
<tr>
<td>mark3jac1200c</td>
<td>8.28</td>
<td>7.87</td>
<td>19.87</td>
</tr>
<tr>
<td>mark3jac1400c</td>
<td>7.85</td>
<td>3.39</td>
<td>4.53</td>
</tr>
<tr>
<td>bayer01</td>
<td>9.56</td>
<td>4.20</td>
<td>5.16</td>
</tr>
<tr>
<td>sinc18</td>
<td>1.58</td>
<td>1.14</td>
<td>0.32</td>
</tr>
<tr>
<td>sinc15</td>
<td>16.54</td>
<td>18.03</td>
<td>22.99</td>
</tr>
<tr>
<td>Zhao2</td>
<td>7.73</td>
<td>3.58</td>
<td>8.77</td>
</tr>
<tr>
<td>multi_dcop.03</td>
<td>2.15</td>
<td>0.90</td>
<td>4.47</td>
</tr>
<tr>
<td>twotone</td>
<td>0.61</td>
<td>0.47</td>
<td>0.24</td>
</tr>
<tr>
<td>onetone1</td>
<td>3.58</td>
<td>2.09</td>
<td>4.69</td>
</tr>
<tr>
<td>tomo1</td>
<td>1.02</td>
<td>0.53</td>
<td>0.86</td>
</tr>
<tr>
<td>bbbmat</td>
<td>44.74</td>
<td>16.50</td>
<td>31.60</td>
</tr>
<tr>
<td>shermanACb</td>
<td>0.26</td>
<td>0.14</td>
<td>0.09</td>
</tr>
</tbody>
</table>

Table 5.4
MA14S ordering, factorization and solution time (in seconds).

Table 5.4 then compares the time of the three main steps of the resolution. Note that the ordering time of both orderings depends on two opposite effects that are
difficult to detect. The better we preserve sparsity, the smaller might be the quotient graph, and the faster we can process it. The better we preserve sparsity, the fewer the elements are absorbed, the fewer the supervariables are detected, and higher the complexity might be. However, our new ordering is a real unsymmetric ordering that selects off-diagonal pivots and updates a constraint matrix. One should thus expect the time spent in the ordering to be higher with CMSL than with DMLS. Indeed, CMSL performs more metric computations and has to explicitly store and manipulate the constraint matrix $C$. The metric update is the most costly step of the ordering so that the complexity of the ordering is tightly linked to the size of $C$ (see Section 4.4.4). Considering that the size of the $C^0$ is typically between $2n$ and $3n$, we see in Table 5.4 that CMSL is quite competitive with respect to DMLS (we observe that the cost of CMSL does not linearly increase with the size of $C^0$). Two algorithmic differences might explain the good behaviour of the CMSL ordering. Since CMSL has the flexibility to select pivots in the constraint matrix it may not be critical to know the metric of the entries that belong to a fairly dense row or column. That is why our CMSL implementation avoids such computation. It sometimes decreases significantly the ordering time (see for example torso1 matrix). Furthermore, supervariables have been generalized in the context to CMSL ordering resulting in separated row and column supervariables. This feature helps CMSL exploit in a more efficient way the unsymmetric structure of the matrix.

We then see in Table 5.3 that the decrease in fill-in and in the number of operations performed during the factorization phase leads to decrease in the factorization time. The gains in factorization time are slightly smaller than the gains in the number of operations (the average decrease in the number of operations is around 21%). This is because sparser factors often leads to smaller full blocks on which basic linear algebra kernels are slower. We observed that the flop rate of MA41JNS tends to be smaller with CMSL than with DMLS: the average flop rate is nearby 1.33 GFlops with the DMLS ordering whereas it is around 1.23 GFlops with the CMSL ordering. The largest gains in time are then obtained for matrices with small flop rate. For example, for the lr* matrices the average flop rate is around 100 MFlops, and MA41JNS is two times faster with CMSL than with DMLS.

5.4. Impact of the hybrid strategies on SuperLU_DIST. Because of the static pivoting strategy used during factorization, SuperLU_DIST is expected to be numerically more sensitive than MA41JNS to the use of hybrid strategies to select pivots during the CMSL ordering. We thus focus in this section, on the analysis of the numerical behaviour of SuperLU_DIST. It has been observed in [3] that, because of static pivoting, iterative refinement may be required to obtain an accurate solution. We thus analyse in this section the component-wise backward error of the solution [6] during iterative refinement. Note that one step of iterative refinement costs at least as much as one forward and backward substitution. The cost of the solution phase is thus very much related to the number of steps of iterative refinement. In the hybrid strategy (see Section 4.1), a relative threshold is set to avoid the selection of small pivots in $C$. The relative threshold was set to 0.01 in all our experiments.

We added to our test set four matrices (see Table 5.5) on which we have observed in [3] that SuperLU_DIST need iterative refinement to improve the accuracy of the solution.

Figure 5.1 compares CMSL and DMLS component-wise backward error during the iterative refinement (results after 0, 2, 3, and 4 steps). From the results it is clear that using the CMSL ordering clearly improves the numerical behavior of SuperLU_DIST.
There are still two matrices (av41092 and Zhan2) for which, after both CMLS and DMLS orderings, iterative refinement does not converge to an accurate solution. The torsol matrix is the only one on which a DMLS approach succeeds whereas a CMLS approach fails. There are seven matrices (sinc18, lhr34c, lhr71c, mult_dcop3, MIXING-TANK, fidap11 and cavity16) on which the backward error of SuperLU_DIST combined with DMLS remains larger than 10^{-10} whereas SuperLU_DIST combined with CMLS converges in less than 4 iterations. Note finally that, with CMLS, on all problems (except torsol, av41092 and Zhan2), three steps of iterative refinement are sufficient to obtain a backward error smaller than 10^{-8}. 

Table 5.5: Test matrices.
6. Concluding remarks. The originality of the CMLS algorithm relies on its ability to compute an unsymmetric permutation with the following goals in mind: to reduce the fill-in in the factors and to preselect numerically good pivots for the factorization. It is based on a constraint matrix which contains the candidate pivots and a quotient graph that is used to compute the structural metrics. The CMLS algorithm can be used to design a family of orderings that can address a large class of problems. The main properties of the algorithm are summarized as follows:

- Significant gains in terms of fill-in (13%) and flops (21%) have been obtained with the structural strategy to select the pivots.
- Using a structural metrics to select the pivots does not affect the numerical behaviour of the MA41JNS solver.
- On numerically difficult problems, CMLS can be used to improve the accuracy of SuperLU_DIST and reduce the number of steps of iterative refinement during the solution phase.
- Although the complexity of the ordering is higher with CMLS than with DMLS, we have shown that CMLS benefits from algorithmic improvements (some of which could even be implemented in DMLS).

The last property emphasizes that DMLS could also benefit from the algorithms developed in the more general and complex framework used for the CMLS algorithm. Both our approach to scale metrics in order to achieve better tie-breaking and our generalized supervariables could be used in the context of DMLS to also improve the metric computation.

One indirect but important consequence of our work is that we do not need to limit our pivot choice to a maximum weighted transversal of the original matrix. Preliminary experiments have shown that the maximum weighted matching can in fact be substituted by a simpler structural maximum transversal during the preprocessing phase (Step 1 as defined in Section 2). One possible direction for future work could then be to design a simpler or easier to parallelize version of the preprocessing phase.

Furthermore, the constraint matrix C contains information that can be seen as an incomplete factorization. We intend to use it as a preconditioner and to compare its quality and cost with existing incomplete LU factorizations.

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